

Métodos Formais  
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## Introduction to Alloy

Área de Teoria DCC/UFMG

# Outline

- Introduction to basic Alloy constructs using a simple example of a static model
  - How to define *sets, subsets, relations with multiplicity constraints*
  - How to use Alloy's *quantifiers* and *predicate* forms
  
- Basic use of the Alloy Analyzer 6
  - *Loading, running, and analyzing* a simple Alloy specification
  - Adjusting basic *tool parameters*
  - Using the *visualization* tool to view instances of models

# Why was Alloy created?

- *Lightweight*
  - Small and easy to use
  - capable of expressing common properties tersely and naturally
  
- *Precise*
  - having a simple and uniform mathematical semantics
  
- *Tractable*
  - amenable to efficient and fully automatic semantic analysis
    - within scope limits

# What is Alloy used for?

- A textual modeling language aimed at expressing *structural* and *behavioral* properties of software systems
- Not meant for modeling code architecture
- But an Alloy specification can be closely related to an OO implementation

# Example applications

- The Alloy distribution comes with several sample models to illustrate the use of Alloy's constructs
  
- Examples
  - Specification of a distributed spanning tree
  - Model of a generic file system
  - Model of a generic file synchronizer
  - Tower of Hanoi model
  - ...

# In summary

- Alloy is general enough that it can model
  - any domain of individuals
  - relations between them
  
- We will start with a few simple examples
  - Not necessarily about software

# Example: Family structure

We want to:

- Model *parent/child relationships* as primitive relations
- Model *spousal relationships* as primitive relations
- Model relationships such as *siblings* as derived relations
- Enforce *biological constraints* via first-order predicates (e.g., people are not their own parents)
- Enforce *social constraints* via first-order predicates (e.g., a spouse isn't a sibling)
- Confirm or refute the existence of certain *derived relationships* (e.g., no one has a spouse with whom they share a parent)

# Atoms and Relations

- In Alloy, everything is built from *atoms* and *relations*
- An *atom* is a primitive entity that is
  - *indivisible*: it cannot be broken down into smaller parts
  - *immutable*: its properties do not change over time
  - *uninterpreted*: it does not have any built in property (the way numbers do for example)
- A *relation* is a structure that *relates atoms*. It is a set of *tuples*, each tuple being a sequence of atoms



# Atoms and Relations: Examples

An *address book* for an email client with a mapping from *names* to *addresses*

FriendBook
Ted -> ted@gmail.com
Ryan -> ryan@hotmail.com

WorkBook
Pilard -> lpilard@ufmg.br
Ryan -> ryan@ufmb.br

- *Unary relations*: a set of names, a set of addresses and a set of books

Name = {(N0), (N1), (N2)}

Addr = {(D0), (D1)}

Book = {(B0), (B1)}

- A *binary relation* from names to addresses

address = {(N0,D0),(N1,D1)}

- A *ternary relation* from books to names to addresses

address = {(B0,N0,D0),(B0,N1,D1),(B1,N1,D2)}

# Relations

- *Size* of a relation: the number of tuples in the relation
- *Arity* of a relation: the number of atoms in each tuple of the relation
  - relations with arity 1, 2, and 3 are said to be unary, binary, and ternary relations

- Examples.

- relation of arity 1 and size 1:

$\text{myName} = \{(N0)\}$

- relation of arity 2 and size 3:

$\text{address} = \{(N0,D0),(N1,D1),(N2,D1)\}$

# Main components of Alloy models

- Signatures and Fields
- Predicates and Functions
- Facts
- Assertions
- Commands and scopes

# Signatures and Fields

- Signatures
  - Describe classes of entities we want to reason about
  - Sets defined in signatures are fixed (dynamic aspects can be modeled by time-dependent relations)
- Fields
  - Define relations between signatures
- Simple constraints
  - Multiplicities on signatures
  - Multiplicities on relations

# Signatures

- A signature introduces a set of atoms
- The declaration

**sig** A {}

introduces a set named A

- A set can be introduced as an extension of another; thus

**sig** A1 **extends** A {}

introduces a set A1 that is a *subset* of A

# Signatures

```
sig A {}  
sig B {}  
sig A1 extends A {}  
sig A2 extends A {}
```

- A1 and A2 are *extensions* of A
- A signature declared independently of any other one is a *top-level signature*, e.g., A and B
- Extensions of the same signature are *mutually disjoint*, as are top-level signatures

# Signatures

```
abstract sig A {}  
sig B {}  
sig A1 extends A {}  
sig A2 extends A {}
```

- A signature can be introduced as a *subset* of another

```
sig A3 in A {}  
sig A2 extends A {}
```

- An *abstract signature* has no elements except those belonging to its extensions or subsets
- All extensions of an abstract signature  $A$  form a *partition* of  $A$

# Fields

- *Relations* are declared as *fields* of signatures
- Writing

**sig** A { f : e }

introduces a relation f of type  $A \times e$ , where e is an expression denoting a product of signatures

- Examples: (with signatures A, B, C)
  - Binary relation:

**sig** A { f1 : B }

where f1 is a subset of  $A \times B$

- Ternary relation:

**sig** A { f2 : B  $\rightarrow$  C }

where f2 is a subset of  $A \times B \times C$



# Example signatures and fields

A family structure:

```
abstract sig Person {  
    children: Person ,  
    siblings: Person  
}
```

```
sig Man, Woman, Other extends Person {}
```

```
sig Married in Person {  
    spouse: Married  
}
```

# Example: family structure

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sig Man extends Person {}  
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sig Other extends Person {}  
sig Married in Person {}
```

An example of an instance is

Person = {(P0), (P1)}

Man = {(P0), (P1)}

Married = {}

Woman = {}

Other = {}

siblings = {(P0,P1), (P1,P0)}

- *siblings* is a binary relation, i.e., a subset of Person x Person
- In the instance, P0 and P1 are siblings

# run Command

- Used to ask AA to generate an instance of the model
- May include *conditions*
  - Used to guide AA to pick model instances with certain characteristics
  - E.g., force certain *sets and relations* to be non-empty
  - In this case, not part of the “true” specification
    - Specific for that run
- We can use conditions to encode *realism constraints* to e.g.,
  - Force generated models to include at least one married person, or one married man, etc.

# run Command

- To analyze a model, you add a run command and instruct AA to execute it.
  - the run command tells the tool to search for an instance of the model
  - you may also give a scope to signatures bounds the size of instances that will be considered
- The scope:
  - *Limits the size of* instances considered to make instance finding feasible
  - Represents the maximum number of elements in a *top-level signature*
  - *Default* value is 3 for each top-level signature
- AA executes only the first run command in a file

## run Example

— The simplest **run** command

— The scope of every signature is 3

```
run {}
```

— The scope scope of every signature is 5

```
run {} for 5
```

— With conditions forcing each **set** to be populated

— Setting the scope to 2

```
run {some Man && some Woman && some Married} for 2
```

— Other scenarios

```
run {some Woman && no Man} for 7
```

```
run {some Man && some Married && no Woman}
```

# Multiplicities

- Allow us to constrain the sizes of sets
  - A multiplicity keyword placed before a signature declaration constrains the number of elements in the signature

`m sig A {}`

- We can also make multiplicities constraints on fields:

`sig A {f: m e}`

`sig A {f: e1 m -> n e2}`

- There are four multiplicities
  - **set** : any number
  - **some** : one or more
  - **lone** : zero or one
  - **one** : exactly one



# Multiplicities: Examples

- Without multiplicity:
  - A set of colors, each of which is red, yellow or green abstract

```
sig Color {}
```

```
sig Red, Yellow, Green extends Color {}
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- A set of colors, each of which is red, yellow or green abstract

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```

- With multiplicity:

- An enumeration of colors

```
abstract sig Color {}  
one sig Red , Yellow , Green extends Color {}
```

# Multiplicities: Examples

- A file system in which each directory contains any number of objects, and each alias points to exactly one object

```
abstract sig Object {}  
sig Directory extends Object {contents: set Object}  
sig File extends Object {}  
sig Alias in File {to: one Object}
```

- The default multiplicity for fields is **one**, so:

```
sig A {f: e}  
sig A {f: one e}
```

are equivalent

# Multiplicities: Examples

- A book maps names to addresses
  - There is at most one address per Name
  - An address is associated to at least one name

```
sig Name, Addr {}  
sig Book {  
    addr: Name some -> lone Addr  
}
```

## Multiplicities: Examples

- A collection of weather forecasts, each of which has a field *weather* associating every city with exactly one weather condition

```
sig Forecast {weather: City -> one Weather}  
sig City {}  
abstract sig Weather {}  
one sig Rainy, Sunny, Cloudy extends Weather {}
```

- Instance:

```
City = {(BH), (Uberlandia)}  
Rainy = {(rainy)}  
Sunny = {(sunny)}  
Cloudy = {(cloudy)}  
Forecast = {(f1), (f2)}  
weather = { (f1, BH, rainy), (f1, Uberlandia, rainy),  
            (f2, BH, rainy), (f2, Uberlandia, sunny) }
```

# Multiplicities and Binary Relations

**sig**  $S$  { **f**: **lone**  $T$  }

- says that, for each element  $s$  of  $S$ ,  $f$  maps  $s$  to *at most* a single value in  $T$

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**sig**  $S \{f: \mathbf{one} T\}$

- Defines a total function

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- For *each element*  $s$  of  $S$ , over the triples that start with  $s$ :  $f$  maps *at most one*  $T$ -element to the same  $V$ -element

# Multiplicities and Relations

- Other kinds of relational structures can be specified using multiplicities

- Examples

- **sig** S {f: **some** T} ... total relation
- **sig** S {f: **set** T} ... partial relation
- **sig** S {f: T **set**  $\rightarrow$  **set** V}
- **sig** S {f: T **one**  $\rightarrow$  V}
- ...

# Cardinality constraints

- Multiplicities can also be applied to whole expressions denoting relations

- `some e`     `e` is non-empty
- `no e`        `e` is empty
- `1one e`      `e` has at most one tuple
- `one e`        `e` has exactly one tuple

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- How would you use multiplicities to define the children relation?



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sig Person {children: set Person}
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## Example: family structure

- How would you use multiplicities to define the children relation?

**sig** Person {children: **set** Person}

- Intuition: each person has zero or more children

- How would you use multiplicities to define the spouse relation?

**sig** Married {spouse: **one** Married}

- Intuition: each married person has exactly one spouse

# Size Determination

- Size determined in a signature declaration has priority on size determined in scope

- Example:

```
abstract sig Color {}  
one sig red, yellow, green extends color {}  
sig Pixel {color: one Color}  
  
run {} for 2
```

- The above limits the signature Pixel to 2 elements, but assigns a size of exactly 3 to Color

# Model weaknesses

- The model is underconstrained
  - It doesn't match our domain knowledge
    - Asymmetric marriage, self child/sibling, asymmetric siblings, multiple fathers...
  - We can add constraints to enrich the model
  
- Under-constrained models are common early on in the development process
  - AA gives us quick feedback on weaknesses in our model
  - We can incrementally add constraints until we are satisfied with it

# Adding constraints

- We'd like to enforce the following constraints (concerning *biology*)
  - No person can be their own parent (or more generally, their own ancestor)
  - No person can have more than one father or mother
  - A person's siblings are those with the same parents
  
- We could also enforce the following *social* constraints
  - The spouse relation is symmetric
  - A man's wife cannot be one of his siblings

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